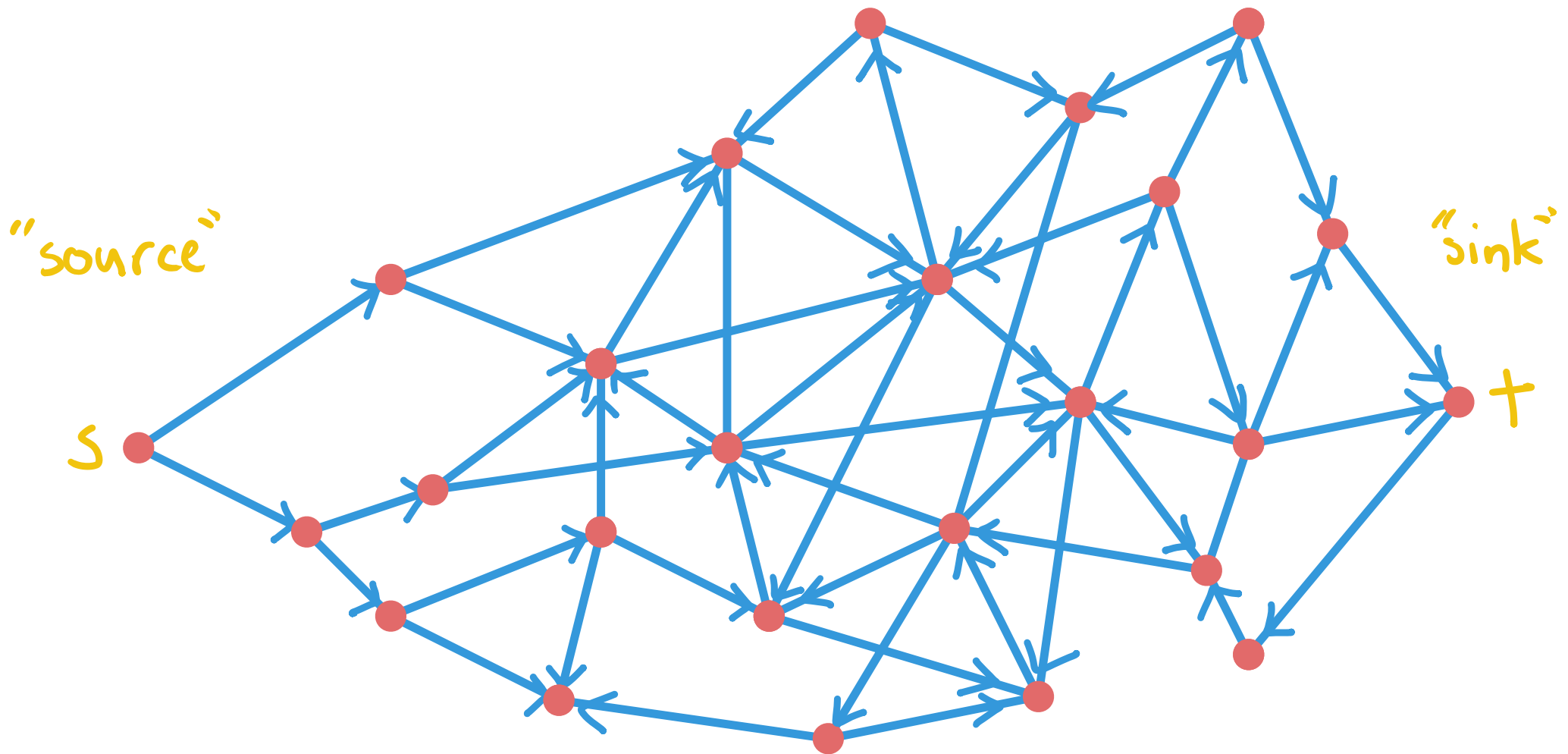
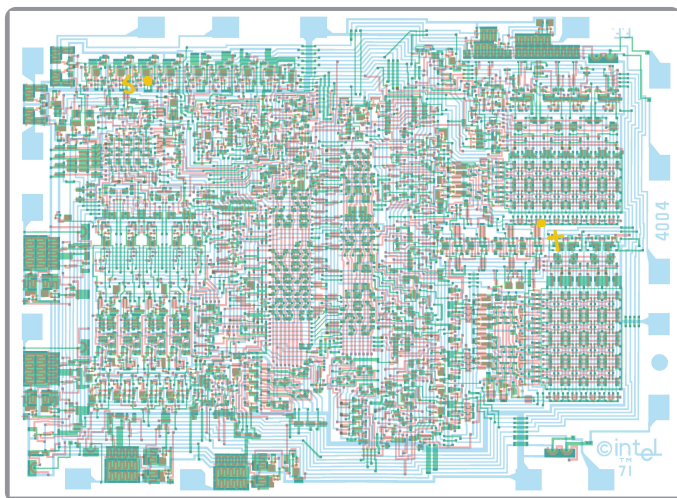
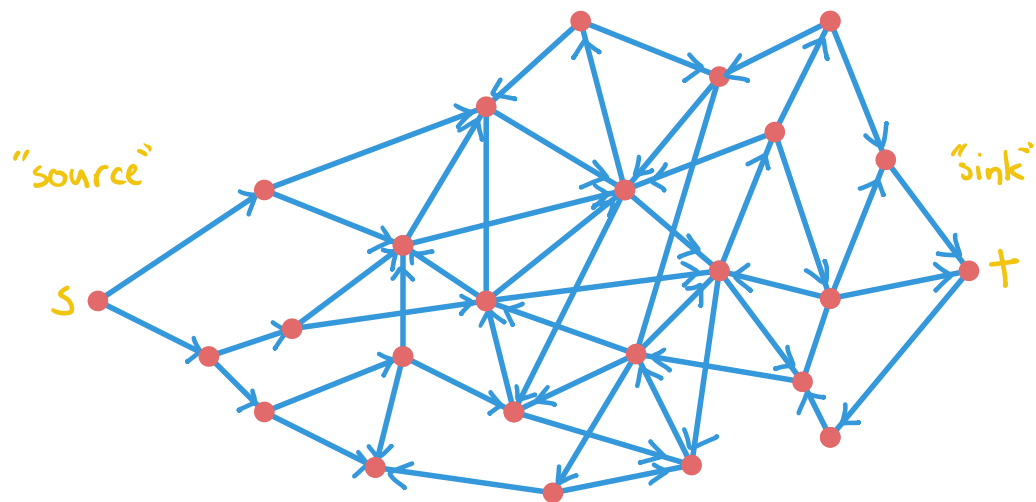


Network Flow

Measuring connectivity

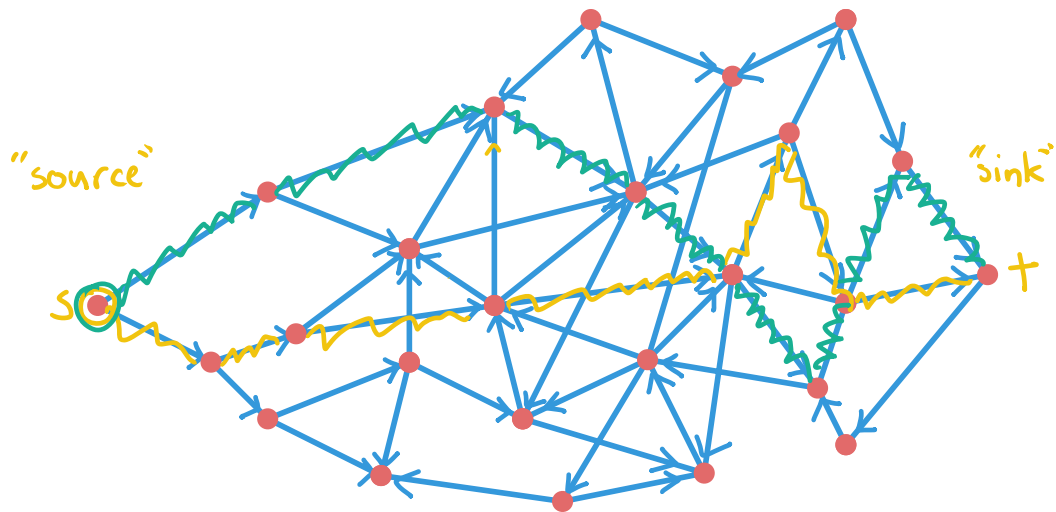


how well connected are s and t?

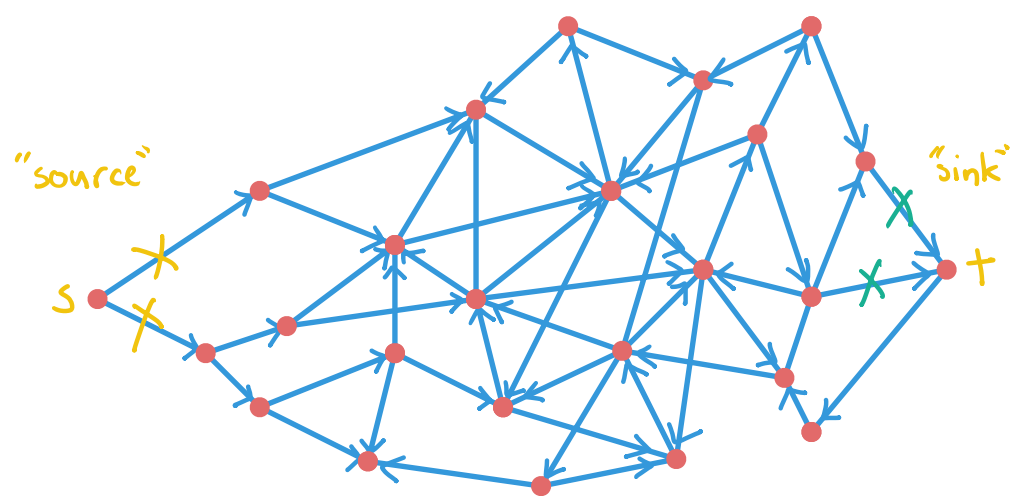


- ① Is there an $s \rightsquigarrow t$ path?
- ② How many distinct $s \rightsquigarrow t$ paths?
- ③ how many k -edge $s \rightsquigarrow t$ walks?
(k fixed)
- ④ length of shortest $s \rightsquigarrow t$ path?
- ⑤ how many shortest $s \rightsquigarrow t$ paths?
- ⑥ In a random walk starting from s , what is the expected # of steps until we reach t ?
- ⑦ max # of edge-disjoint $s \rightsquigarrow t$ paths?
- ⑧ min # of edge deletions to disconnect $s \rightsquigarrow t$?

Measuring connectivity

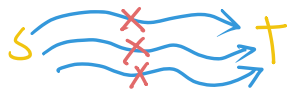


⑦ max # of edge-disjoint $s \rightarrow t$ paths?



⑧ min # of edge deletions to disconnect $s \rightarrow t$?

path $p \xrightarrow{\text{packs into}} \text{edge } e \in P$
 $\xleftarrow{\text{Covers}}$



path $p \xrightarrow{\text{packs into}} \text{edge } e \in P$
 $\xleftarrow{\text{Covers}}$

Lemma

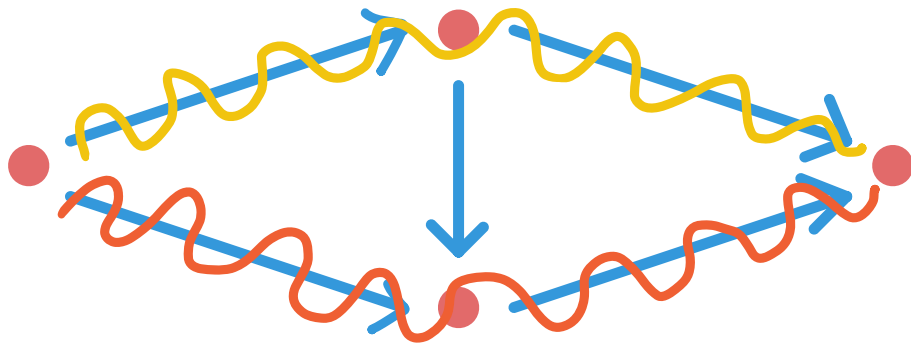
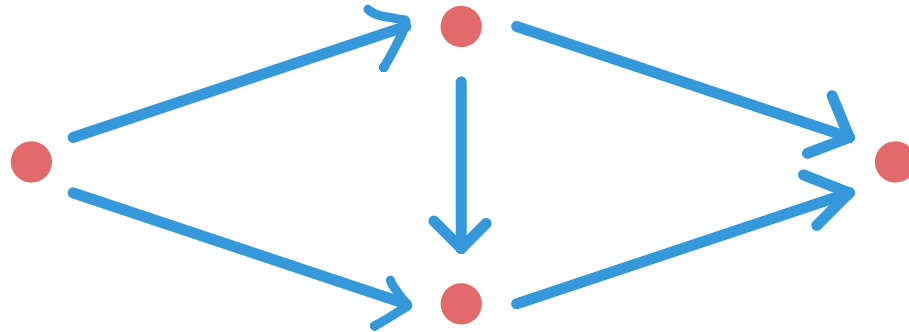
$$\left(\begin{array}{c} \text{max \#} \\ \text{disjoint} \\ s \rightsquigarrow t \text{ paths} \end{array} \right) \leq \left(\begin{array}{c} \text{min \# edges} \\ \text{in any } s \rightsquigarrow t \text{ cut} \end{array} \right)$$

* 2nd example of
packing $\leq$ covering

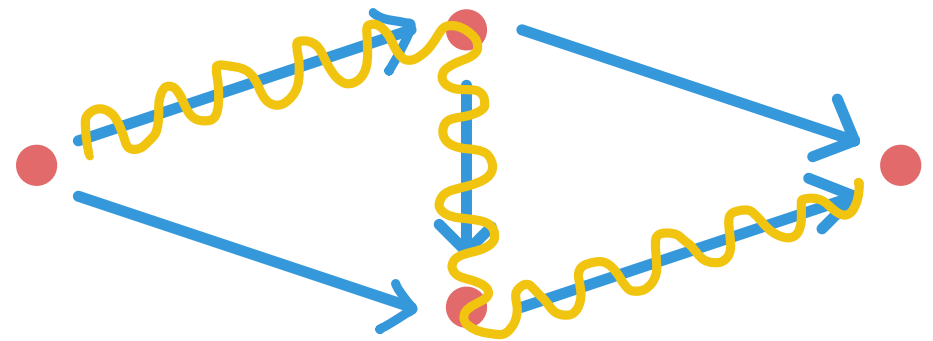
Iterative/greedy

① while there is a path $p: s \rightsquigarrow t$

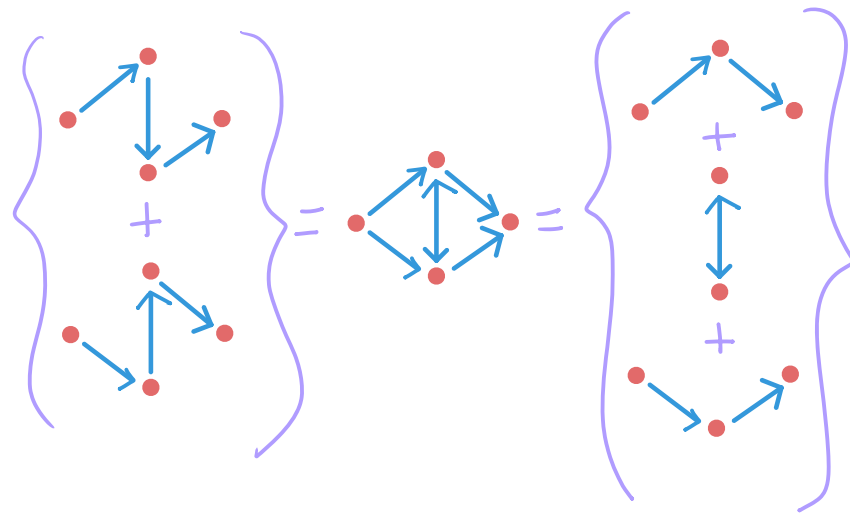
@ take p and remove all its edges from the graph



(good)



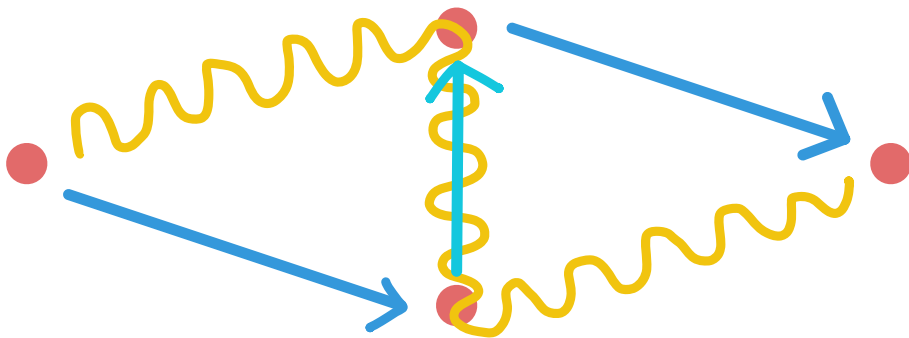
(bad)



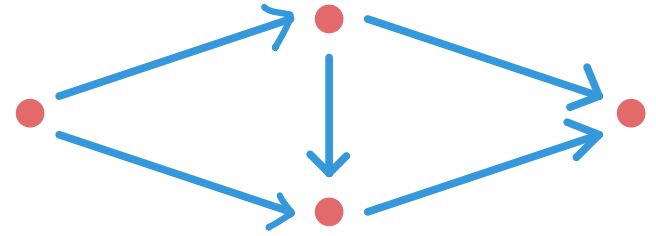
"opposite edges"



"cancel out"



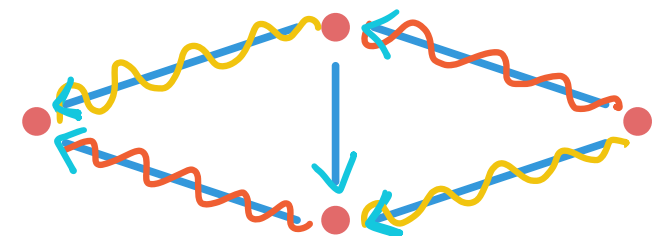
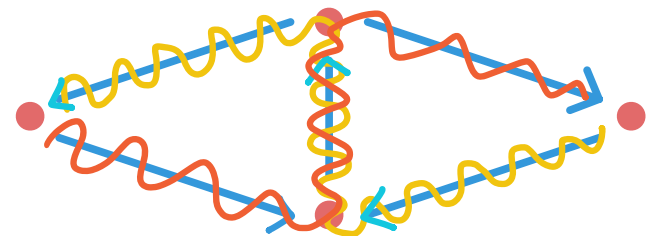
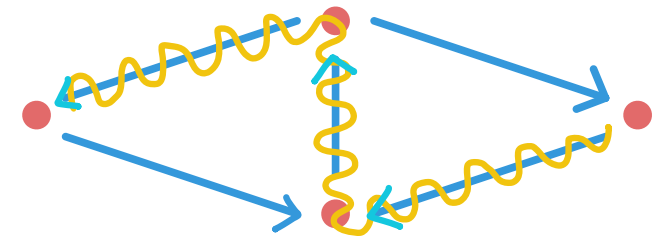
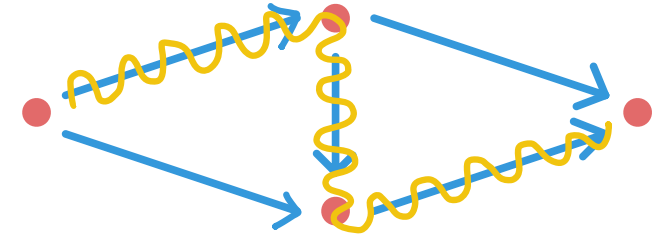
we should be allowed
to take reverse edges
in subsequent paths



augmenting-paths($G = (V, E), s, t$)

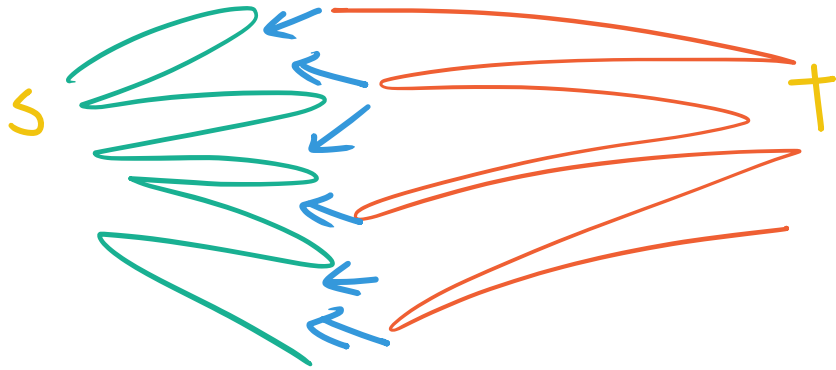
/ Ford and Fulkerson [FF56]*

1. $P \leftarrow \emptyset$
2. While there is an (s, t) -path $p : s \rightsquigarrow t$:
 - A. Remove all of p 's edges from the graph.
 - B. Add the reverse of all of p 's edges to the graph.
 - C. Uncross p with previously selected paths, and add p to P .



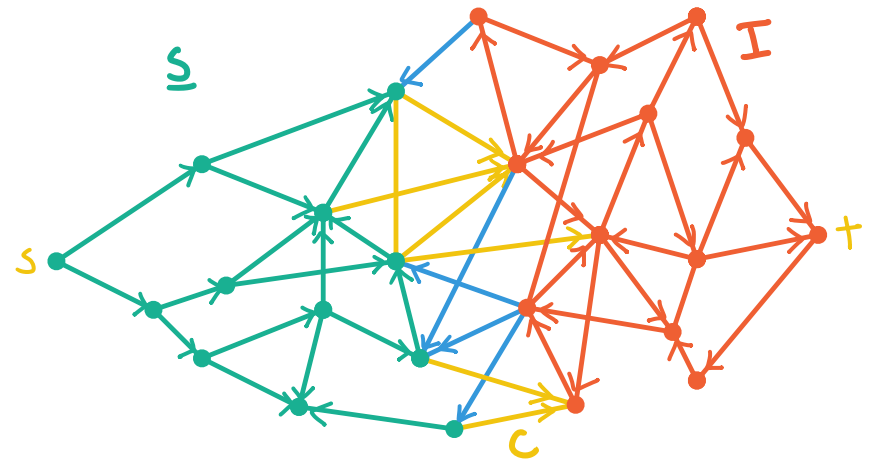
Proving optimality

algorithm terminates w/ $s \rightarrow t$ disconnected
in residual graph



let $\underline{S}$ = vertices reachable from s in residual graph
 $\underline{I} = V \setminus \underline{S}$

$C = \delta^+(\underline{S})$ = edges from $\underline{S}$ to $\underline{I}$



over the course of the algorithm:

- ① each $s \rightarrow t$ path has (exactly) 1 more edge from $\underline{S}$ to $\underline{I}$ than from $\underline{I}$ to $\underline{S}$
- ② each $s \rightarrow t$ path decreases $|C|$ by 1
- ③ At the end, $|C| = 0$.

$$\Rightarrow (\text{\# paths packed}) = (\text{\# edges in } C)$$

$$(\text{\# paths packed}) = (\text{max \# disjoint } s \rightarrow t \text{ paths}) = (\text{min \# edges in any } s \rightarrow t \text{ cut}) = (\text{\# edges in } C)$$

equal

$$(\# \text{ paths packed}) \stackrel{=}{=} \left(\begin{array}{c} \text{max \#} \\ \text{disjoint} \\ s \rightarrow t \text{ paths} \end{array} \right) \stackrel{=}{=} \left(\begin{array}{c} \text{min \# edges} \\ \text{in any } s \rightarrow t \text{ cut} \end{array} \right) \stackrel{=}{=} \left(\begin{array}{c} \text{\# edges} \\ \text{in } C \end{array} \right)$$

Theorems

$$\textcircled{1} \left(\begin{array}{c} \text{max \#} \\ \text{disjoint} \\ s \rightarrow t \text{ paths} \end{array} \right) = \left(\begin{array}{c} \text{min \# edges} \\ \text{in any } s \rightarrow t \text{ cut} \end{array} \right)$$

$$\textcircled{2} (\# \text{ paths packed}) = \left(\begin{array}{c} \text{max \#} \\ \text{disjoint} \\ s \rightarrow t \text{ paths} \end{array} \right)$$

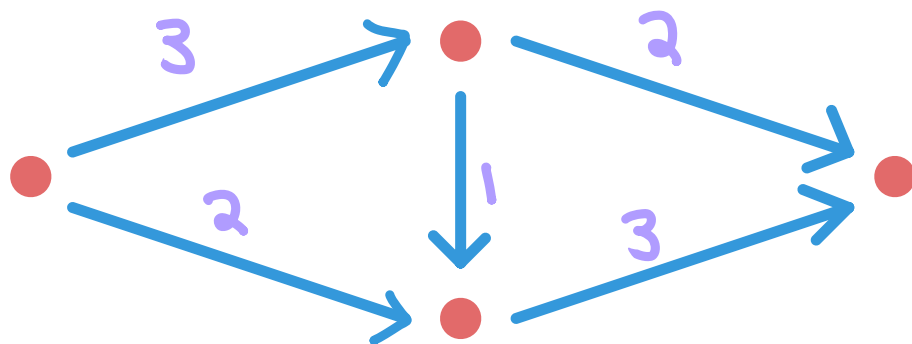
$$\textcircled{3} \left(\begin{array}{c} \text{min \# edges} \\ \text{in any } s \rightarrow t \text{ cut} \end{array} \right) = \left(\begin{array}{c} \text{\# edges} \\ \text{in } C \end{array} \right)$$

1. $P \leftarrow \emptyset$
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 - C. Uncross p with previously selected paths, and add p to P .³
3. $S \leftarrow \{\text{vertices still reachable from } s\}$.
4. Return the (s, t) -path packing P and the (s, t) -cut $\partial^+(S)$.

*/

Theorem 14.5 ([FF56]). Let $G = (V, E)$ be a directed graph with m edges, and let $s, t \in V$ be distinct vertices. Let λ be the size of the max-flow/min-cut from s to t . Then $\text{augmenting-paths}(G, s, t)$ produces both a maximum packing and a minimum cut in $O(m\lambda)$ time.

Capacities



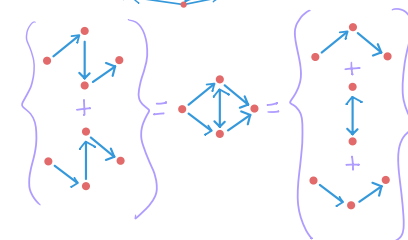
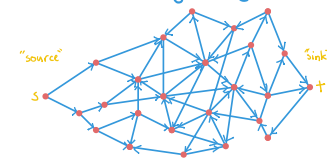
capacity $c(e) > 0$: reuse edge e up to $c(e)$ times

- if $c(e)$ large, Augmenting-Paths not polynomial time

- if $c(e)$ not integral (and just real-valued)
what about max fractional path packing

How to compute

⑦ max # of edge-disjoint $s \rightsquigarrow t$ paths?



Lemma let p, q be edge-disjoint $s \rightsquigarrow t$ paths



then there is p', q' : $s \rightsquigarrow t$ disjoint without opposite pairs of edges



$$\textcircled{1} \left(\begin{array}{l} \text{max \#} \\ \text{disjoint} \\ s \rightsquigarrow t \text{ paths} \end{array} \right) = \left(\begin{array}{l} \text{min \#} \\ \text{edges in any} \\ s \rightsquigarrow t \text{ cut} \end{array} \right)$$

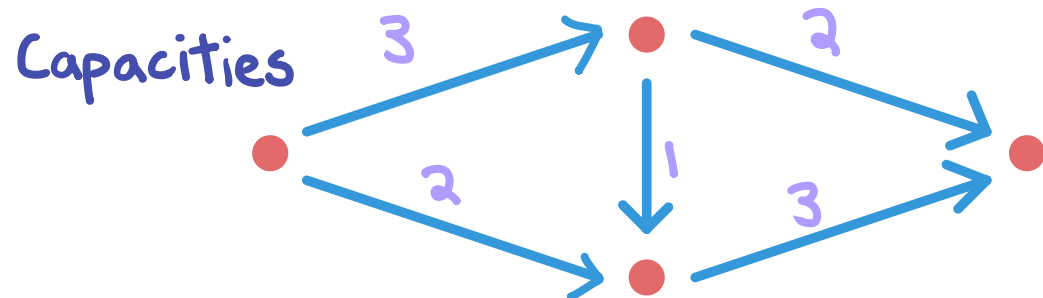
$$\textcircled{2} \left(\begin{array}{l} \text{\# paths} \\ \text{packed} \end{array} \right) = \left(\begin{array}{l} \text{max \#} \\ \text{disjoint} \\ s \rightsquigarrow t \text{ paths} \end{array} \right)$$

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*/

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capacity $c(e) > 0$: reuse edge e up to $c(e)$ times

• if $c(e)$ large, Augmenting-Paths not polynomial time

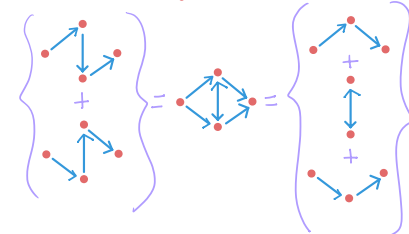
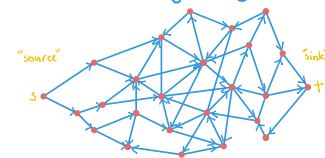
• if $c(e)$ not integral (and just real-valued) what about max fractional path packing

Good news:

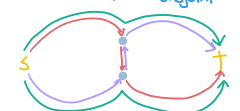
Our duality theorems still hold

How to compute

⑦ max # of edge-disjoint $s \rightsquigarrow t$ paths?



Lemma let p, q be edge-disjoint $s \rightsquigarrow t$ paths



then there is $p', q': s \rightsquigarrow t$ disjoint without opposite pairs of edges



$$\textcircled{1} \left(\begin{array}{c} \text{max \#} \\ \text{disjoint} \\ s \rightsquigarrow t \text{ paths} \end{array} \right) = \left(\begin{array}{c} \text{min \#} \\ \text{edges in any} \\ s \rightsquigarrow t \text{ cut} \end{array} \right)$$

$$\textcircled{2} \left(\begin{array}{c} \text{\# paths} \\ \text{packed} \end{array} \right) = \left(\begin{array}{c} \text{max \#} \\ \text{disjoint} \\ s \rightsquigarrow t \text{ paths} \end{array} \right)$$

$$\textcircled{3} \left(\begin{array}{c} \text{min \#} \\ \text{edges in any} \\ s \rightsquigarrow t \text{ cut} \end{array} \right) = \left(\begin{array}{c} \text{\# edges} \\ \text{initially in } C \end{array} \right)$$

remaining issues algorithmic ~ next time ~

Max Flow "liquid" version of path packing

Input $G=(V,E)$ directed capacities $c:E \rightarrow \mathbb{R}_{>0}$

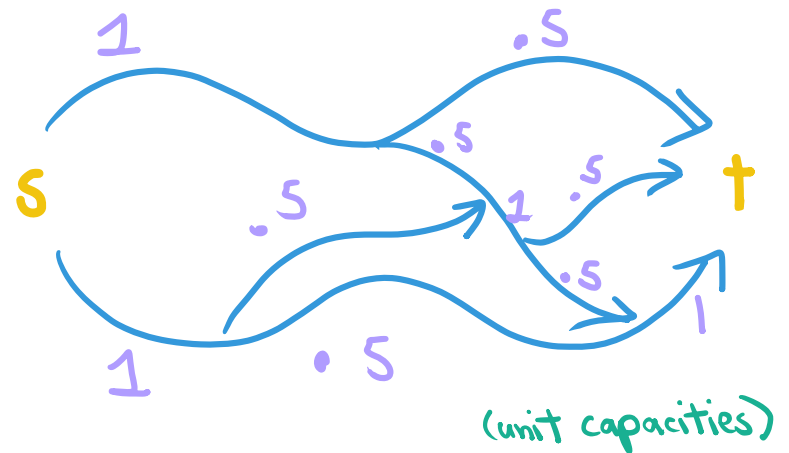
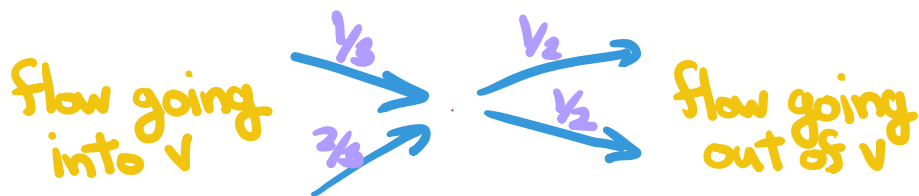
source $s \in V$, sink $t \in V$

"s-t flow" $f:E \rightarrow \mathbb{R}_{\geq 0}$ s.t.

(a) "capacity constraints"
 $0 \leq f(e) \leq c(e) \quad \forall e \in E$

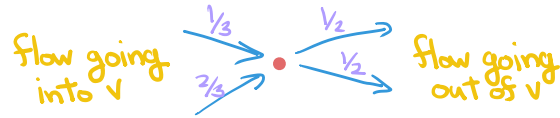
(b) "conservation of flow"

$$\sum_{(u,v) \in E} f(u,v) = \sum_{(v,w) \in E} f(v,w) \quad \forall v \notin \{s,t\}$$

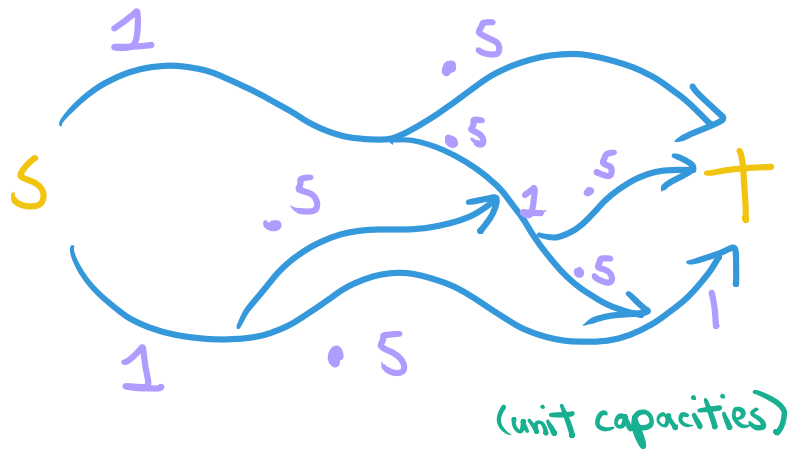


(b) "Conservation of flow"

$$\sum_{(u,v) \in E} f(u,v) = \sum_{(v,w) \in E} f(v,w) \quad \forall v \notin \{s,t\}$$



net flow out of s = net flow into t



"net flow from s to t",
"size of flow"

Max flow problem

find $s \rightarrow t$ flow of maximum size

"linear program"

$$\max \underbrace{\sum_{(s,w) \in E} f(s,w) - \sum_{(u,s) \in E} f(u,s)}_{\text{size of flow}} \quad \text{over } f \in \mathbb{R}^E$$

"capacity constraints"

$$\text{s.t. } 0 \leq f(e) \leq c(e) \quad \forall e \in E$$

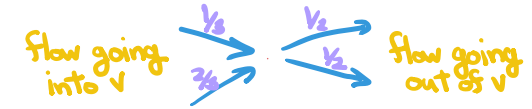
"conservation of flow"

$$\sum_{(u,v) \in E} f(u,v) = \sum_{(v,w) \in E} f(v,w) \quad \forall v \notin \{s,t\}$$

"s-t flow" $f: E \rightarrow \mathbb{R}_{\geq 0}$ s.t.

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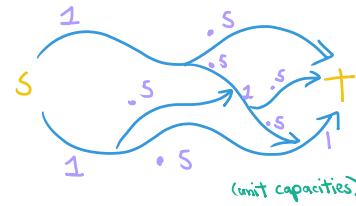
(b) "conservation of flow"
 $\sum_{(u,v) \in E} f(u,v) = \sum_{(v,w) \in E} f(v,w) \quad \forall v \notin \{s,t\}$



"linear program" because objective, constraints, all linear functions of f

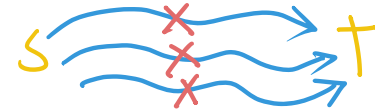
How does

Max Flow

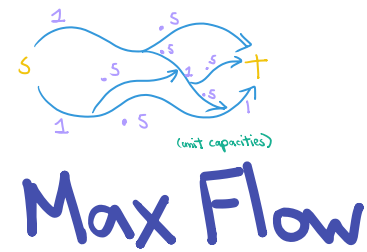


relate to

max path packings, min cut?



 max path packings vs



let $\chi: \mathcal{P}_{s,t} \rightarrow \mathbb{R}_{\geq 0}$ encode a path packing

define $f: E \rightarrow \mathbb{R}_{\geq 0}$ as:

$$f(e) = \sum_{p: e \in p} \chi_p.$$

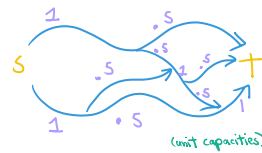
then f is feasible flow of same size

(Exercise about converse)



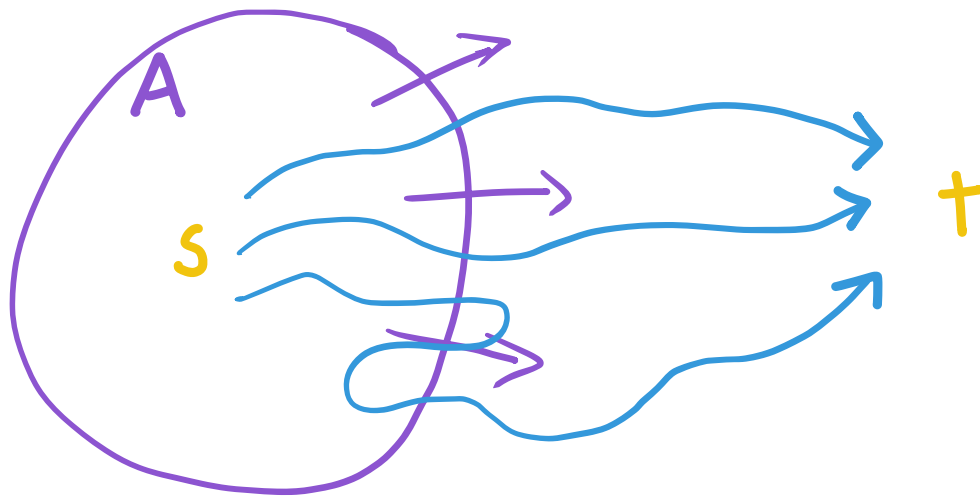
min cut

vs



Max Flow

let $C = \mathcal{J}^+(A)$ min-cut, $f: E \rightarrow \mathbb{R}_{\geq 0}$ max flow
 ($A \subset V$ w/ $s \in A, t \notin A$)



$$(\text{size of } f) = (\text{net flow out of } A) \leq (\text{capacity of } \mathcal{J}^+(A))$$

max path packings vs Max Flow

let $x: P_{s,t} \rightarrow \mathbb{R}_{\geq 0}$ encode a path packing

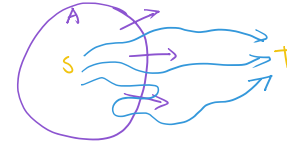
define $f: E \rightarrow \mathbb{R}_{\geq 0}$ as:

$$f(e) = \sum_{p: e \in p} x_p$$

then f is feasible flow of same size

min cut vs Max Flow

let $C = \partial^+(A)$ min-cut, $f: E \rightarrow \mathbb{R}_{\geq 0}$ max flow
($A \subset V$ w/ $s \in A, t \notin A$)

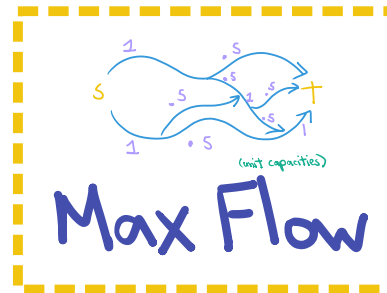


$$(\text{size of } f) = (\text{net flow out of } A) \leq (\text{capacity of } \partial^+(A))$$

we showed



$\leq$



$\leq$



equal

max path packings vs MaxFlow

let $x: P_{st} \rightarrow \mathbb{R}_{\geq 0}$ encode a path packing

define $f: E \rightarrow \mathbb{R}_{\geq 0}$ as:

$$f(e) = \sum_{p \ni e} x_p$$

then f is feasible flow of same size

min cut vs MaxFlow

let $C = \mathcal{J}^c(A)$ min-cut, $f: E \rightarrow \mathbb{R}_{\geq 0}$ max flow

(MFC of unit flow)

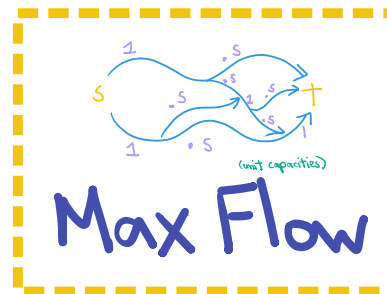


(size of f) = (net flow out of A) $\leq$ (capacity of $\mathcal{J}^c(A)$)

Theorem



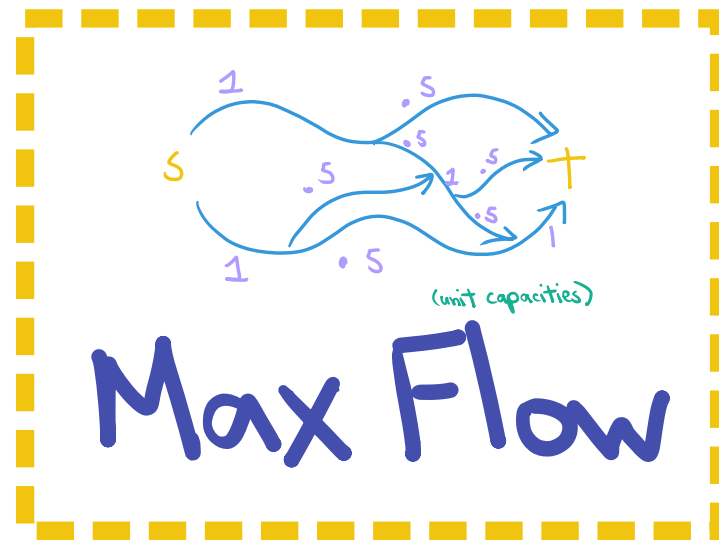
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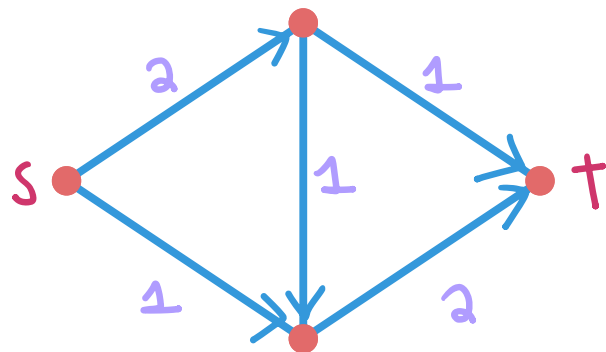
Next: strongly polynomial times



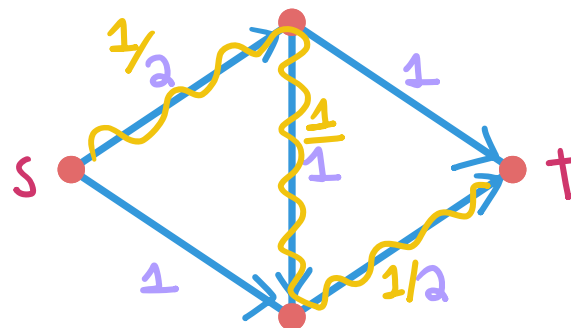
Residual graphs w/r/t flows

(exactly what you would expect)

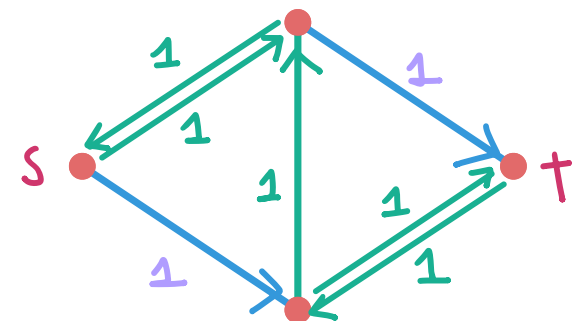
(capacities)



(1 unit flow)

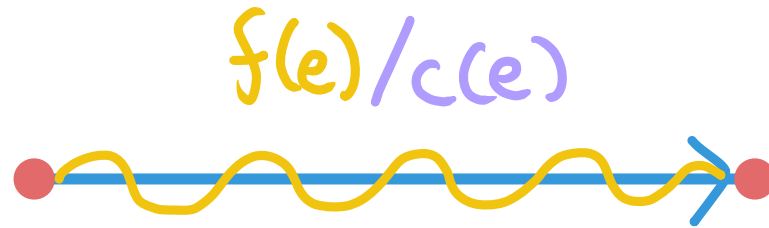


(residual edges)



let $f: E \rightarrow \mathbb{R}_{\geq 0}$ be (s,t) -flow

$e \in E$ w/ capacity $c(e)$

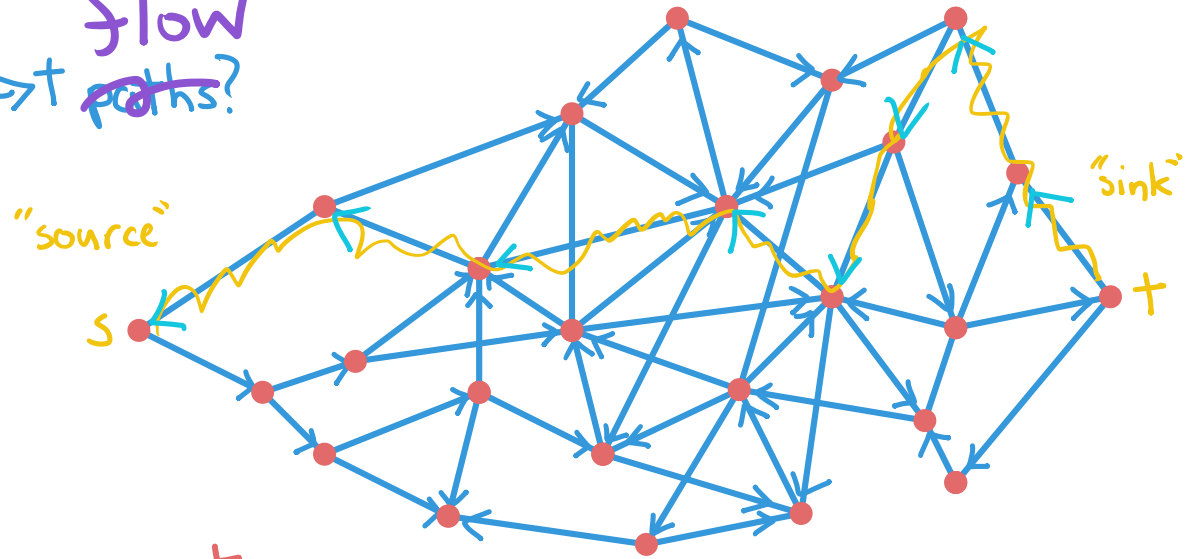


this leaves $f(e) - c(e)$ capacity in forwards direction
and $f(e)$ capacity backwards



How to compute

⑦ max # of ~~edge-disjoint~~ $s \rightsquigarrow t$ ~~paths?~~ ^{flow}



① while there is a path $p: s \rightsquigarrow t$

Ⓐ take p and ~~remove all its edges from the graph~~

route as much flow as possible

Ⓑ add the reverse of all of p 's edges in the graph

Ⓒ uncross p w/ previously selected paths

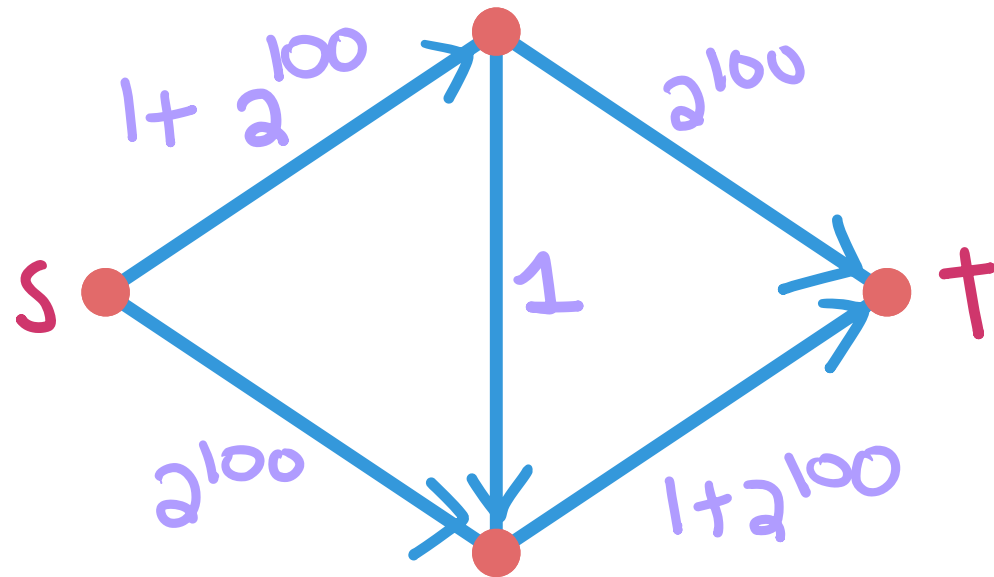
update residual graph

↪ reasonable, but not quite polynomial time

① while there is a path $p: s \rightsquigarrow t$:

② take p and route as much flow as possible

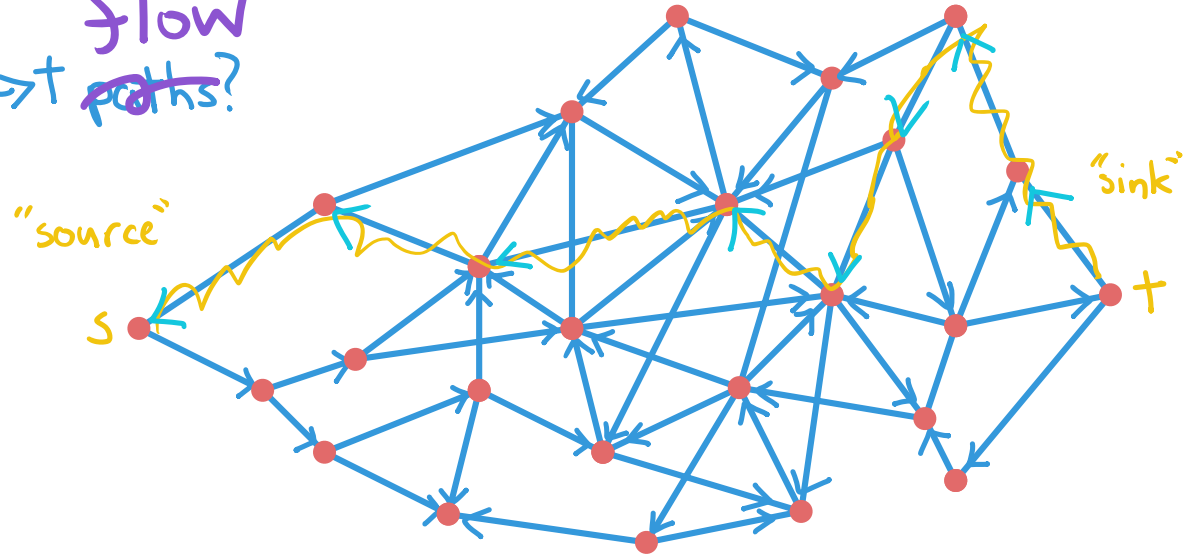
③ update residual graph



How to compute

⑦ max # of ~~edge-disjoint~~ $s \rightsquigarrow t$ ~~paths?~~ ^{flow}

shortest path
(Edmonds, Karp)



① while there is a path $p: s \rightsquigarrow t$

Ⓐ take p and ~~remove all its edges from the graph~~

route as much flow as possible

Ⓑ add the reverse of all of p 's edges in the graph

Ⓒ uncross p w/ previously selected paths

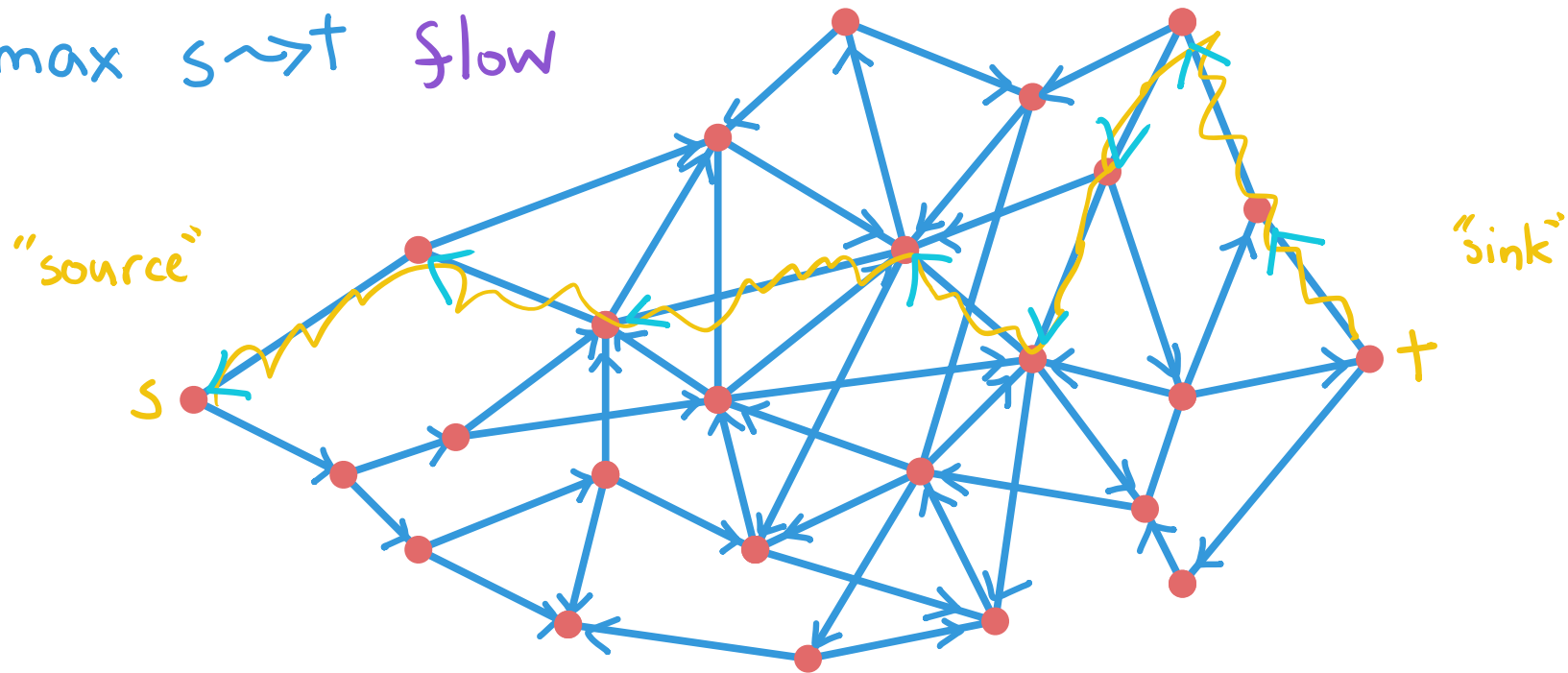
update residual graph

(cleaned up)

(Edmonds, Karp)

How to compute

⑦ max $s \rightsquigarrow t$ flow



① while there is a path $p: s \rightsquigarrow t$:

Ⓐ take shortest path p and route as much flow as possible

Ⓑ update residual graph

Analysis

- ① while there is a path $p: s \rightsquigarrow t$:
 - Ⓐ take shortest path p and route as much flow as possible
 - Ⓑ update residual graph

Key
Lemma

$\forall v \in V$, distance from s to v in residual graph never decreases

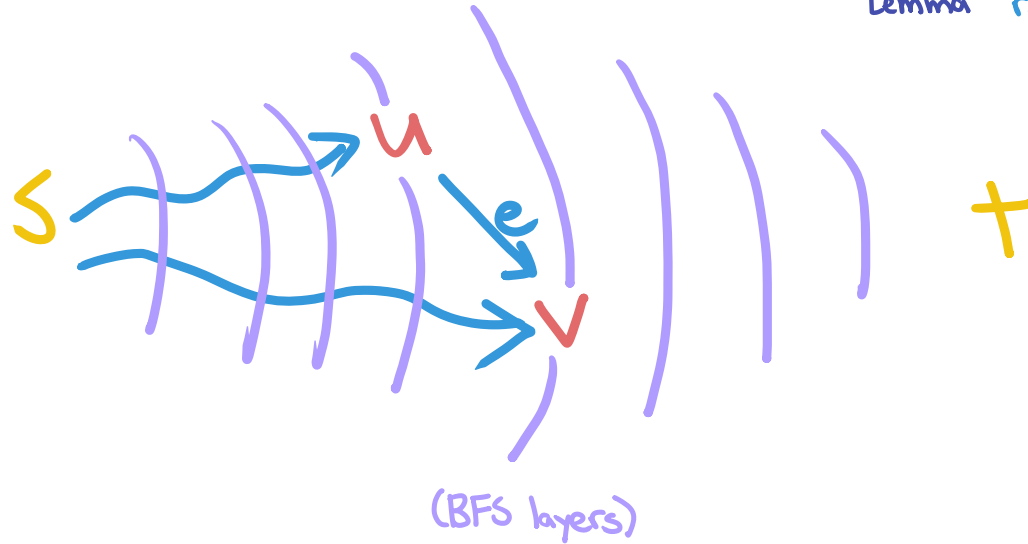
let us first assume lemma is true

fix $e=(u \rightarrow v) \in E$.

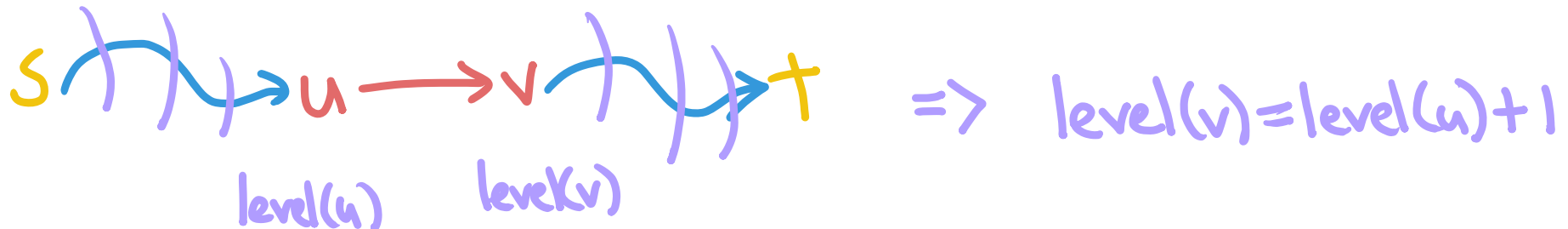
consider distances from s to u & v

- ① while there is a path $p: s \rightsquigarrow t$:
 - Ⓐ take shortest path p and route as much flow as possible
 - Ⓑ update residual graph

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lemma residual graph never decreases



① e removed from ^(residual graph) RG , $\Rightarrow e$ in shortest $s \rightsquigarrow t$ path



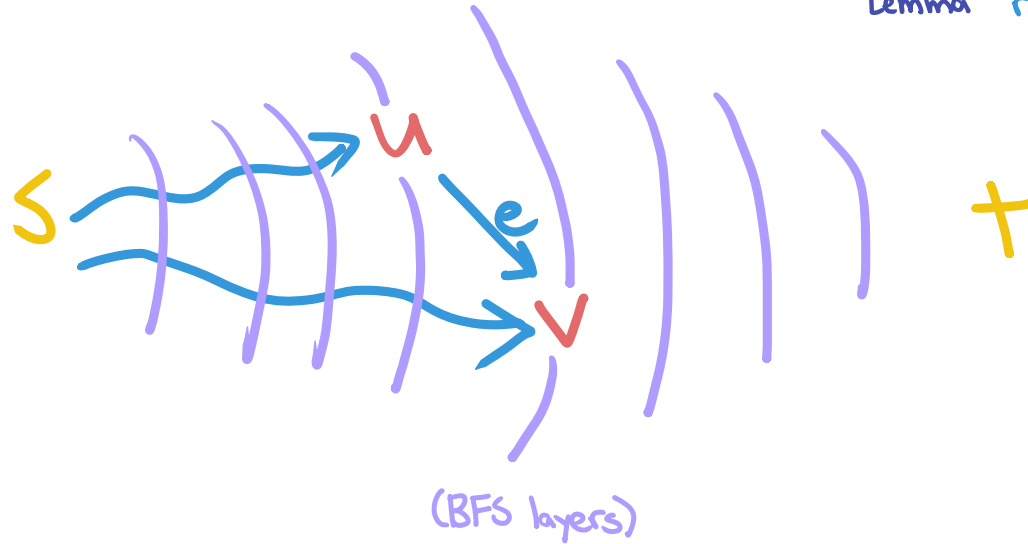
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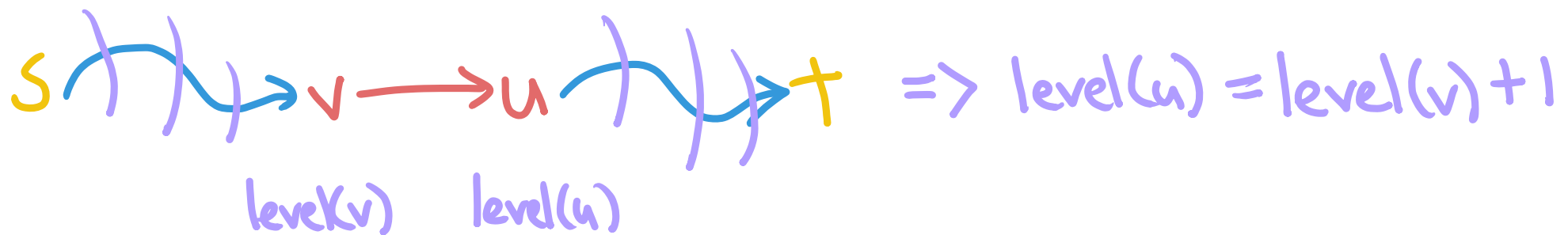
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② $e = (u \rightarrow v)$ added back in to RG

$\Rightarrow (v \rightarrow u)$ was in shortest $s \rightsquigarrow t$ path



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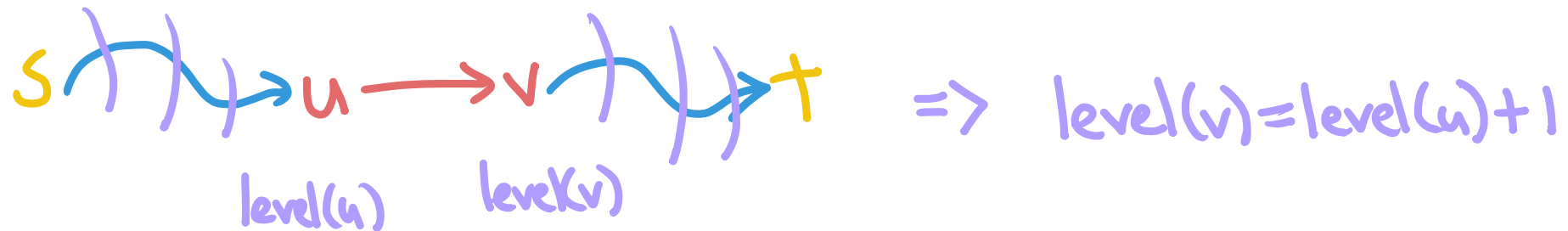
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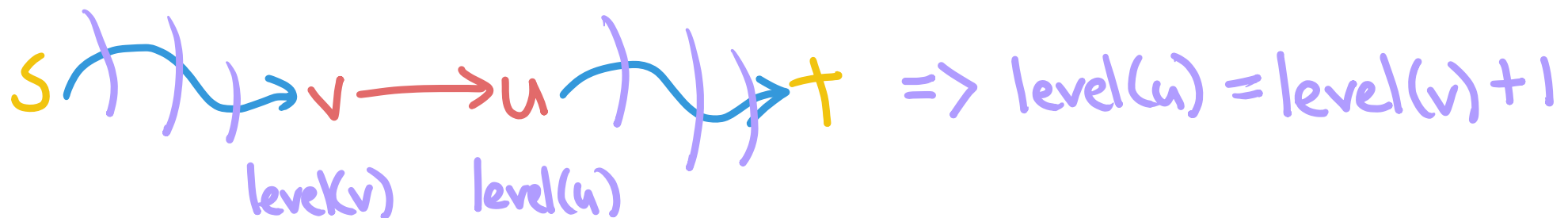
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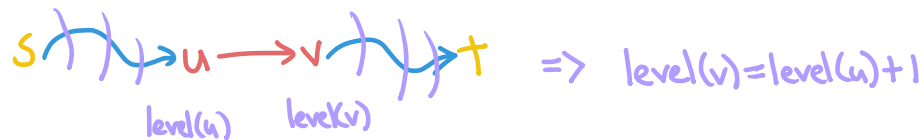
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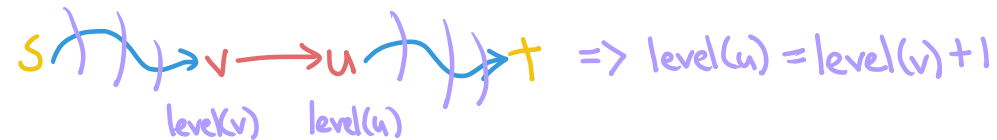
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each event ①/② $\Rightarrow \text{level}(u), \text{level}(v)$ have switched order.

levels nondecreasing by key lemma

max finite level is $n-1$

$\Rightarrow e$ gets deleted/added at most n times

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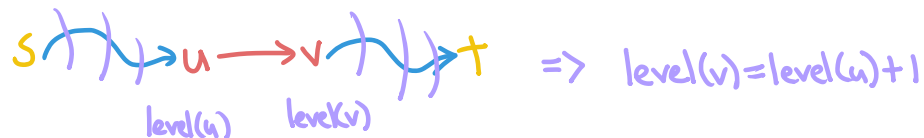
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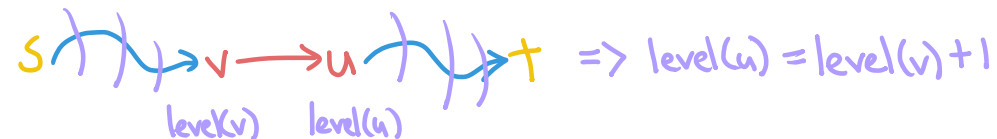
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✓ e gets deleted/added at most n times

Last observation each iteration removes (at least)
one edge from residual graph. why?

b/c we route as much as possible

= $O(mn)$ iterations

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Running time

$O(m)$ per iteration

$O(mn)$ iterations

$O(m^2n)$ time

Key $\forall v \in V$, distance from s to v (" $l(v)$ ")
Lemma in residual graph never decreases

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proof

suppose we augment along $p: s \rightsquigarrow t$
need to show no levels decrease

proof by induction on level after updating residual graph.

Base case: level 0

$s \rightsquigarrow \rightsquigarrow \rightsquigarrow$

s is the only vertex at level 0, which never changes

General case: level $\beta > 0$

let $l(v) = d$ before updating
 $l(v) = \beta > 0$ after updating

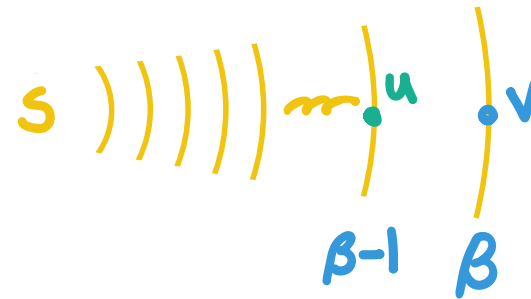
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Before



After



$u =$ 2nd to last vertex on shortest (s, v) path after

level(u) = $\beta-1$ after augmenting

$\Rightarrow$ level(u) $\leq \beta-2$ before augmenting

$\Rightarrow (v \rightarrow u)$ not on augmenting path

$\Rightarrow (u \rightarrow v)$ in residual graph
before augmenting

$\Rightarrow \text{level}(v) \leq \text{level}(u) + 1 \leq \beta - 1$
before augmenting

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Edmonds, Karp

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computes max (s,t) -flow, min (s,t) -cut in $O(m^2n)$ time

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